

# Finite-Repetition threshold for infinite ternary words

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## Overview

- Fraenkel and Simpson result,
- Dejean's Conjecture on the Repetitive threshold,
- Overlap-free words and number of squares they contain,
- Finite-Repetition Threshold, Shallit's Theorem,
- Finite-Repetition Threshold, main result
- Conclusion
- Generalizing the result for larger alphabet sizes?

# Infinite Binary with 3 squares

## Repetitions

Example:  $\underbrace{abaab}_{\text{period}} abaabab$

Length=12

Period=abaab

Exponent =  $\frac{\text{length}}{\text{periodlength}} = \frac{12}{5} = 2.4$

Square: exponent even integers.

e.g abab

Cube: exponent integers multiple of 3.

e.g aaa

## Fraenkel and Simpson Theorem

There is an infinite Binary word containing 3 squares 00, 11 and 0101 only.

## Example

011000111001011100111000101110010110001011100011001...

- Fraenkel, Simpson 1995 (complicated construction)

- Harju, Nowotka 2006

$$\begin{cases} 0 \rightarrow 111000110010110001110010 \\ 1 \rightarrow 111000101100011100101100010 \\ 2 \rightarrow 111000110010110001011100101100 \end{cases}$$

- Ochem 2006 (50-uniform)

$$\begin{cases} 0 \rightarrow 00011001011000111001011001110001011100101100010111 \\ 1 \rightarrow 00011001011000101110010110011100010110001110010111 \\ 2 \rightarrow 00011001011000101110010110001110010111000101100111 \end{cases}$$

- Badkobeh, Crochemore 2010

$$0 \rightarrow 012, 1 \rightarrow 02, 2 \rightarrow 1.$$

$$\begin{cases} 0 \rightarrow 01001110001101 \\ 1 \rightarrow 0011 \\ 2 \rightarrow 000111 \end{cases}$$

## Thue-Morse morphism $t$

$$B = \{0, 1\}$$

$$\begin{cases} t(0) & = 01, \\ t(1) & = 10. \end{cases}$$

$$t^\infty(0) = 0110100110010110\dots$$

$t^\infty(0)$  is overlap-free.

Overlap is  $2^+$ -repetition in the form  $auaua$  e.g.  $\underbrace{abc}_{abc} a$ .

## Squares are unavoidable

# Avoidability on 3 letters

The morphism  $f_0$  is defined from  $A$  to itself

$$A = \{a, b, c\}$$

$$\begin{cases} f_0(a) &= abc, \\ f_0(b) &= ac, \\ f_0(c) &= b. \end{cases}$$

## Iteration and Properties

$$f_0(a) = abc,$$

$$f_0^2(a) = abcacb,$$

.

.

.

$$f_0^\infty(a) = abcacbabcabacabcacb \dots$$

It is known that  $f_0^\infty(a)$  is square-free.

# Avoidability on 3 letters

## The Dejean's morphism on 3 letters

The Dejean's morphism  $d$  is defined from  $A$  to itself

$$\begin{cases} d(a) &= \text{abcacbcabcbacbcacba}, \\ d(b) &= \text{bcabacabcacbacabacb}, \\ d(c) &= \text{cabcbabcabacbabcbac}. \end{cases}$$

$d^\infty(a) = \text{abcacbcabcbacbcacbabcabacabcacbacabacb} \dots$

The  $d^\infty(a)$  is  $(\frac{7}{4})^+$ -free.

## $\frac{7}{4}$ power is unavoidable

The longest word on 3 letters that can avoid  $(\frac{7}{4})$  power has length 39.

# Repetitive Threshold

- $\text{MaxExp}(x) = \max\{ e \in \mathbb{Q} \mid u^e \text{ occurs in } x \text{ for some } u \}$
- Threshold,  $t_k$  for  $k$ -letter alphabet is  $t$  if
  - Exists infinite string  $x$ ,  $\text{MaxExp}(x) = t$
  - no infinite string  $y$ , satisfies  $\text{MaxExp}(y) < t$
- Example  $t_2 = 2$ ,  $t_3 = \frac{7}{4}$



# Dejean's Conjecture on the Repetitive threshold

## Dejean's Conjecture

The repetition threshold for a  $k$ -letter alphabet (with  $k \geq 2$ ) is:

$$\begin{cases} 2 & \text{for } k = 2 \\ \frac{7}{4} & \text{for } k = 3 \\ \frac{7}{5} & \text{for } k = 4 \\ \frac{k}{k-1} & \text{otherwise.} \end{cases}$$

## Proofs

- *J.J. Pansiot*. A Propos d'une Conjecture de F. Dejean sur les Répétitions dans les Mots. (1984)
- *J. Moulin Ollagnier*. Proof of Dejean's Conjecture for alphabets with 5, 6, 7, 8, 9, 10 and 11 letters. (1992),
- *A. Carpi*. On Dejean's conjecture over large alphabets. (2007)
- *J. Currie, N. Rampersad*. Dejean's Conjecture holds for  $n \geq 27$  (2009)
- *M. Rao*. Last Cases of Dejean's Conjecture.  $8 \leq k \leq 34$ . (2009).
- *J. Currie, N. Rampersad*. Dejean's Conjecture holds for  $n \geq 30$  (2009)

# Overlap-free Binary words and number of squares they contain

## Observations

Maximum length of an overlap-free binary word with at most  $s$  squares.

$s$	= 0	1	2	3	4	5	6	7	8	9	10	11
$l(s)$	= 3	5	8	12	14	18	24	26	32	38	50	54
$s$	= 12	13	14	15	16	17	18	19	20	21	22	23
$l(s)$	= 66	78	102	110	134	158	206	222	270	318	414	446

## Theorem

[Shallit's theorem (2008)] *For all  $t \geq 1$ , there are no infinite binary words that simultaneously avoid all squares  $yy$  with  $|y| \geq t$  and  $\frac{7}{3}$  powers.*

Consequently there is no infinite overlap-free binary word with a bounded number of squares.

## Finite-Repetition Threshold

Finite-Repetition Threshold on  $k$ -letter alphabet,  $FRt(k)$ , is  $f$  if

- Exists an infinite  $k$ -letter word  $x$  with  $\text{MaxExp}(x) = f$  and with finitely many  $t_k$ -powers. ( $t_k$  is Dejean's threshold)
- No infinite  $k$ -letter word  $y$  with  $\text{MaxExp}(y) < f$  contains finitely many  $t_k$ -powers.

Example:  $k = 2$ ,  $f = \frac{7}{3}$

$FRt(2) = \frac{7}{3}$ , Badkobeh and Crochemore (JM2010)

*There exists an infinite binary word whose factors have an exponent at most  $7/3$  and that contains **12 squares**, the fewest possible.*

12 is the minimum number.

Maximum length of  $(7/3)^+$  free binary words with at most  $s$  squares.

$s$	= 0	1	2	3	4	5	6	7	8	9	10	11
$\ell(s)$	= 3	5	8	12	14	18	24	30	37	43	83	116

# Main Result

- The *finite-repetition threshold* of the 3-letter alphabet is its Dejean's repetition threshold, that is,  $7/4$ .
- The smallest number of  $7/4$ -powers occurring in a  $7/4^+$ -power free infinite ternary word is 2, the minimum number.
- The largest  $7/4^+$ -power free ternary words with only one  $7/4$ -power is 102 letter long.

The two unavoidable  $7/4$ -powers occurring in the word are  $(0121)^{7/4} = 0121012$  and  $(2010)^{7/4} = 2010201$ .  
Up to a permutation of letters.

# Morphism

$g$  from  $\{a, b, c, d\}^*$  to  $\{0, 1, 2\}^*$  defined by:

$$\left\{ \begin{array}{l} g(a) = 0102101202102010210121020120210120102120121020120210121 \\ 0212010210121020102120121020120210121020102101202102012 \\ 10212010210121020120210120102120121020102101210212, \\ g(b) = 0102101202102010210121020120210120102120121020120210121 \\ 0201021012021020121021201021012102010212012102012021012 \\ 10212010210121020120210120102120121020102101210212, \\ g(c) = 0102101202102010210121020120210120102120121020102101202 \\ 1020121021201021012102010212012102012021012102010210120 \\ 21020102120121020120210120102120121020102101210212, \\ g(d) = 0102101202102010210121020120210120102120121020102101202 \\ 1020102120121020120210121020102101202102012102120102101 \\ 21020102120121020120210120102120121020102101210212. \end{array} \right.$$

The morphism is uniform with codeword length 160. The morphism  $g$  translates any infinite  $7/5^+$ -free word on the alphabet  $a, b, c, d$  into a  $7/4^+$ -free ternary word containing only two  $7/4$ -powers, the fewest possible.

# Proof based on Ochem's result

## Ochem's Lemma

Let  $\alpha, \beta \in \mathbb{Q}$ ,  $1 < \alpha < \beta < 2$ , and  $p \in \mathbb{N}^*$ . Let  $h : \Sigma_s^* \rightarrow \Sigma_e^*$  be a synchronizing  $q$ -uniform morphism (with  $q \geq 1$ ). If  $h(w)$  is  $(\beta^+, p)$ -free for every  $\alpha^+$ -free word  $w$  such that  $|w| < \max\left\{\frac{2\beta}{\beta-\alpha}, \frac{2(q-1)(2\beta-1)}{q(\beta-1)}\right\}$ , then  $h(t)$  is  $(\beta^+, p)$ -free for every (finite or infinite)  $\alpha^+$ -free word  $t$ .

## Application to morphism $g$

$\alpha = 7/5$  and  $q = 160$ .

Choose  $\beta = 17/10$  and  $n = 5$ .

$g(w)$  is  $(17/10^+, 5)$ -power free for all words  $w$  for which

$$|w| < \max\left\{\frac{2\beta}{\beta-\alpha}, \frac{2(q-1)(2\beta-1)}{q(\beta-1)}\right\},$$

which implies  $|w| < 12$ .

# An infinite ternary word with only one square

The morphism  $h$  from  $\{a, b, c, d\}^*$  to  $\{0, 1, 2\}^*$  defined by:

$$\begin{cases} h(a) = 010020102120 \\ h(b) = 010021020120 \\ h(c) = 021001201020 \\ h(d) = 021012021020. \end{cases}$$

The morphism is uniform with codeword length 12.

The morphism  $h$  translates any infinite  $7/5^+$ -free word on the alphabet  $\{a, b, c, d\}$  into a overlap-free ternary word containing only one square, the fewest possible.

# Conjecture and Questions

## Conjecture

The *finite-repetitions threshold* of 4-letter alphabet is  $\frac{7}{5}$  and the associated number of  $\frac{7}{5}$ -powers is 2.

List of Repetitions:  $\{(acdbcadcbd)^{\frac{7}{5}}, (bcdacbdcad)^{\frac{7}{5}}\}$ ,

Computation shows that the maximal length of  $(\frac{7}{5})^+$  free 4-ary word with only one  $\frac{7}{5}$ -repetition is 230.

Conjecture tested up to length 100000.

## Questions

- Finite Repetition Threshold for larger alphabet sizes?